

# Tree induction

\* note how we split the large tree of height  $k$  into smaller subtrees of height  $k-1$  or shorter \*

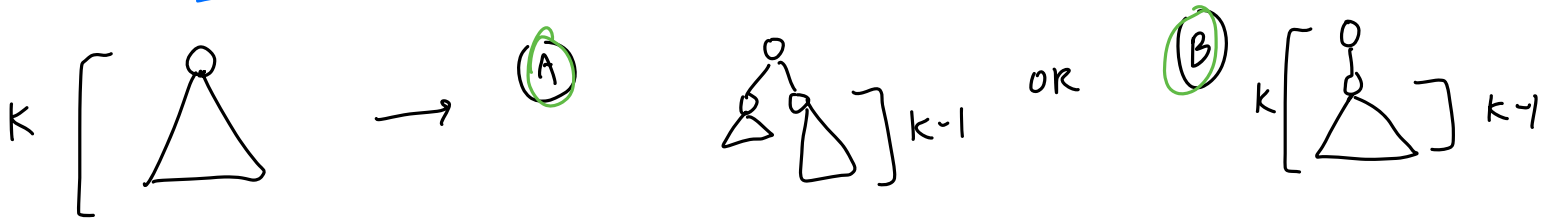
tree example) Let  $T$  be a binary tree with height  $h$  and  $n$ -many nodes.  
Then,  $n \leq 2^{h+1} - 1$ .

Proof by induction on height  $h$ .

Base:  $h=0$   $1 \leq 2^1 - 1 = 1$  ✓  
 $n=1$

IH: Suppose binary tree of height  $h$  &  $n$ -many nodes has  $n \leq 2^{h+1} - 1$  for  $h=1 \dots k-1$ .

Goal: for a tree of height  $k$ , it's true that  $n \leq 2^{k+1} - 1$ .



Inductive step:

case 1: Assume we have a tree of height  $k$ . case 1 is that the root node has 2 children. Then, each of those children is the root node of a tree of height  $\leq k-1$ . The left subtree then has  $\leq 2^k - 1$  nodes,  $\downarrow$  by IH. The right subtree also has  $\leq 2^k - 1$  nodes,  $\downarrow$  by IH.

Then our tree of height  $k$  has  $\leq 1 + 2^k + 2^k - 1$  nodes.  
 $n \leq 2 \cdot 2^k - 1 = 2^{k+1} - 1$

Case 2: Assume we have a tree of height  $k$ . Case 2 is that the root node has 2 child of height  $= k-1$ . By IH, this subtree has  $\leq 2^{k-1}$  nodes. So our tree has  $n \leq 2^{k-1} + 1 = 2^k \leq 2^{k+1} - 1$  since  $k > 0$ .

Grammar example)

consider grammar  $G$  with start  $S$  and terminals  $a, b$ .

$S \rightarrow \underbrace{ab}_{\text{base case}} \mid \underline{SS} \mid \underline{aSb}$ .

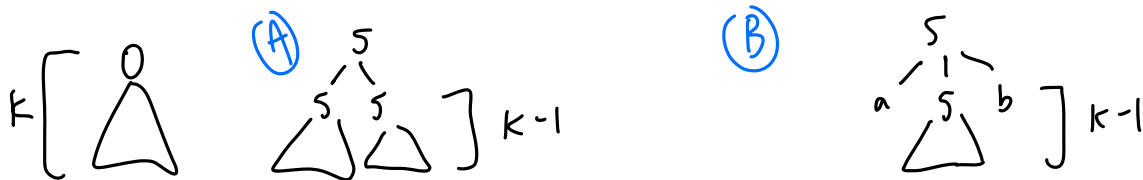
Claim: All trees generated by  $G$  will have the same number of  $a$  nodes and  $b$  nodes.

Proof by induction on height  $h$ .

Base:  ~~$h=0$~~   $S$   $h=1$   $\left[ \begin{array}{c} S \\ / \quad \backslash \\ a \quad b \end{array} \right]$  The number of  $a$ 's and  $b$ 's are equal (1).

IH: Suppose all trees generated by  $G$  of height  $h$  will have the same # of  $a$  and  $b$  nodes for  $h = 1 \dots k-1$ .

Goal: Show tree of height  $k$  generated by  $G$  has equal #  $a$ 's &  $b$ 's.



Case 1: Our tree of height  $k$  has 2 children rooted at  $S$ , call the subtrees  $S_1$  and  $S_2$ . Both  $S_1$  and  $S_2$  have height  $\leq k-1$ . By IH,  $S_1$  has  $m$ -many  $a$ 's and  $b$ 's, and  $S_2$  has  $n$ -many  $a$ 's and  $b$ 's,  $m, n \in \mathbb{N}$ . So, our tree has  $m+n$ -many  $a$ 's and  $m+n$ -many  $b$ 's.

Case 2: Our tree of height  $k$  has 3 children, rooted at  $a S b$ .  
The subtree rooted at  $S$  has  $n$ -many  $a$ 's and  $b$ 's.  
Our tree then has  $n+1$ -many  $a$ 's and  $b$ 's from the left & right children.